Unit 3 Notes: Graphs and Circuits

Section 3.1: Modeling with Graphs

Konigsberg Bridge Problem: "Is it possible for a person to walk around the city crossing each bridge exactly once, starting and ending at the same point?"

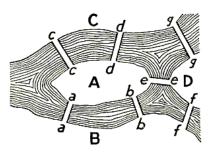
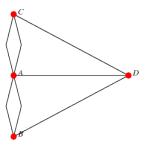


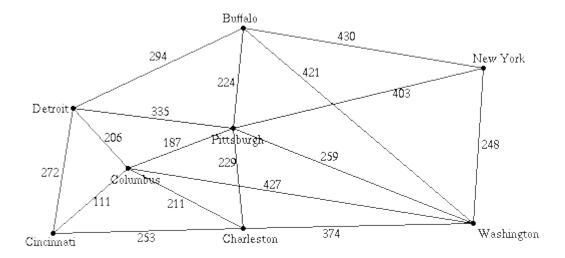
Figure 98. Geographic Map: The Königsberg Bridges.



Euler constructed the simplified geometric model of the situation called a graph, thus creating the subject of **graph theory**. These types of graphs are not the same as graphs of functions or relations (like in Algebra). Graphs are used to solve a wide variety of problems.

ACTIVITY: What's the Shortest Route?

Suppose a salesperson wishes to travel to each city in the map below exactly once, starting and ending in New York, and using only the roads shown. The numbers on the roads indicate distances (in miles) between cities. Find the shortest route that the salesperson could use.



Work in small groups to come up with your route and the justification for why you chose the route you did. Look for a method that will work every time you try a problem like this. Are you convinced that your route is the shortest? Why?

Example 1:

Building a house is usually a team effort that involves specialists, such as architects, excavators, concrete workers, framing carpenters, roofers, etc. Different specialists are often able to work at the same time provided that the work that must precede a particular specialist is completed before that specialist begins. By working simultaneously whenever possible, the house can be completed more quickly, and it is natural to wonder if there is an optimal way to schedule the various tasks for completion.

Information to build a house:

	<u>Task</u>	Time (Days)	Immediately Preceding
			<u>Task</u>
Α	Preparing final house and site plans	3	None
В	Excavation and foundation construction	5	A
С	Framing and closing main structure	12	В
D	Plumbing	5	С
Е	Wiring	3	С
F	Heating-cooling installation	7	E
G	Insulation and dry wall	9	D, F
Н	Exterior siding, trim, and painting	15	С
I	Interior finishing and painting	7	G
J	Carpeting	3	I
Κ	Landscaping	4	Н

If we did each task individually, it would take 73 working days to complete the house. Draw a graph we could use to help the builder decide which tasks can be done simultaneously in order to complete the job more quickly. Then use the graph to determine the least number of days to complete the house.

Notice that the algorithm used in the example above is recursive (depends on the previous term). Here is the general algorithm to calculate the number of days for a particular task:

If there are no prerequisite tasks, use the number of days required by the task alone.

Otherwise:

- (1) Calculate the number of days for each prerequisite task by using the algorithm.
- (2) Choose the largest of the numbers found in step 1, and add to it the number of days required by this task alone.

Definition: When you can travel along each edge of	of a graph in only one directio	n, the graph is
called a	or	•
(A sample of such a graph is a probability tree!)		

Example 2:

60% of the students in a college live on campus. 70% of those who live on campus favor a tuition increase to pay for improved student health services. 40% of those who live off campus favor this increase. Draw a graph to determine the following:

- (a) What is the proportion of students who favor this increase to pay for improved student health services?
- (b) If a randomly selected student favors the tuition increase, what is the probability that the student lives on campus?

Section 3.2: The Definition of a Graph

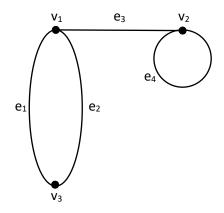
Аε	ra	рh	consist	s of:
----	----	----	---------	-------

1	A finite set of	
Ι.	A finite set of	

2. A finite set of _______,

3.	A function (called the	_ function) that maps each edge to a	
	set of either one or two vertices (the	of the edge).	

Consider the graph and table below:



edge	endpoints
e_1	$\{v_{1}, v_{3}\}$
e ₂	$\{v_{1}, v_{3}\}$
e ₃	$\{v_{1}, v_{2}\}$
e ₄	{v ₂ }

The vertices are:	<u>.</u>	The edges are:	
Definition: This table/graph r	epresents an		function.
The essential feature of an eccurve) is not important!	dge is that it	its endpoints	. The SHAPE (or
Definition: Two vertices conn	nected by an edge are ₋		·
Definition: Two edges with a	common endpoint are	called	
CAUTION: Edges in graphs (un endpoints. Just because two	• • • • • • • • • • • • • • • • • • • •		
Note that v_1 and v_3 are connection	ected by more than one	e edge. When this occu	rs then we say that
the edges are	Also, edge $\it e_{\it 4}$	$_{\scriptscriptstyle \downarrow}$ joins vertex v_2 to itsel	f. Such an edge is
called a Ir	all graphs there MIIS	The a vertex at each e	nd of each edge

Example 1:

Draw a picture of the graph *G* defined as follows.

1. Set of vertices: $\{v_1, v_2, v_3, v_4, v_5\}$

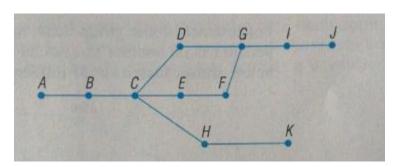
2. Set of edges: $\{e_1, e_2, e_3, e_4, e_5\}$

3. Edge-endpoint function:

Edge	Endpoints
e_1	$\{v_1, v_2\}$
e_2	$\{v_1, v_4\}$
e_3	$\{v_1, v_4\}$
e_4	$\{v_{5}\}$
e_5	$\{v_4v_5\}$

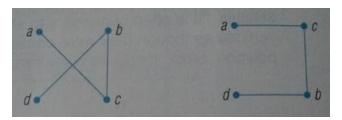
Notice that v_3 is not the endpoint of any edge. That vertex is called an ______. Although all edges must have endpoints, a vertex need not be the endpoint of an edge.

Take a look at the house-building problem from 4.1:



In this graph, there are no loops or parallel edges. Such a graph is called ______.

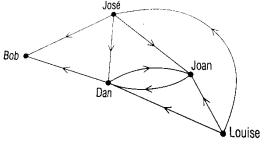
Consider the simple graph with vertices a, b, c, d and edges $\{a$, $c\}$, $\{b$, $d\}$, and $\{b$, $c\}$. Two pictures of this graph are shown below. The picture on the left has two edges that cross, but note that they do NOT intersect (there is not a vertex at the point of intersection). This type of intersection is called a _______. The figure on the right illustrates the same graph, but it avoids all crossings.



Example 2: Draw all simple graphs with vertices $\{u, v, w\}$ if one of the edges is $\{u, v\}$.

Sometimes it is useful to add direction to each edge of a graph. We call this kind of graph a **digraph**. Instead of writing $\{v_1, v_2\}$ for the endpoints of an edge, in a digraph we would write the endpoints as an ordered pair (v_1, v_2) .

For instance, some group-behavior studies investigate the influence one person has on another in a social setting. The directed graph pictured below shows such a set of influence relationships. (NOTE: This is the same thing as the communication networks we learned about in Unit 2!)



Who does Jose influence?

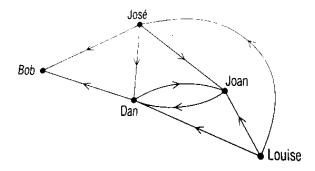
Who does Louise influence?

Surprisingly, we can describe a graph numerically, even when no numbers are given. We can do this using a matrix! (YEAH!)

Definition: The **adjacency matrix** M for a graph with vertices $v_1, v_2, ..., v_n$ is the $n \times n$ matrix in which the element in the ith row and jth column is the number of edges from vertex v_i to vertex v_j .

Example 3:

Write the adjacency matrix for the directed graph of influence relationships pictured below.



Example 4:

Draw a picture of a graph (not directed) that has the following adjacency matrix.

$$\begin{bmatrix} 1 & 2 & 0 \\ 2 & 2 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

Section 3.3: Handshake Problems

In groups of 6, work through the following problems:

In sub-groups of the size indicated, everyone shakes hands with everyone else. Record your data in the table below:

Number of People	Total Number of
in Group	Handshakes
1	
2	
3	
4	
5	
6	

Represent these sets of data with graphs. Label your graphs H_n , where n is the number of people in a group.

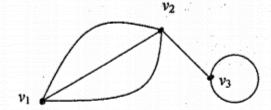
$$H_1$$
 H_2 H_3

$$H_4 \hspace{1cm} H_5 \hspace{1cm} H_6$$

 H_n is called a **complete graph** because every vertex is connected to every other vertex by an edge. Since every person shakes hands once with each other person, every pair of vertices is joined by exactly one edge.

Definitions: If v is a vertex of a graph G, the degree of v, denoted deg(v), equals the number of edges that have v as an endpoint, with each edge that is a loop counted twice. The **total degree of** G is the sum of the degrees of all the vertices of G.

Example - Given the graph, G



$$deg(v_1) = 3$$

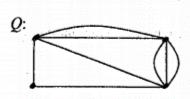
$$deg(v_2) = \underline{\qquad}$$

$$deg(v_3) = \underline{\qquad}$$

Total degree of
$$G =$$

Branch Barry Branch Carlot Carlot

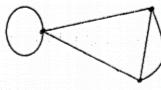
Now examine the following graphs and fill in the chart below.



Graph

Q

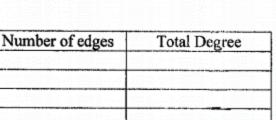
S T H₂ R:



S:



T:



Can you find a relationship between the number of edges in a graph a	and the total degree?

9

Why do you think this relationship exists?

What does this relationship tell you about the total degree of any graph?

Now pick 5 people in your group. Have each of the 5 people shake hands with exactly 3 others in the group of 5. What happens?
Try to draw a graph of this situation.
Now try to have the group of 5 shake hands with exactly 1, 2, and 4 others. When does this experiment fail? When is it successful?
How about a different sized group? Try some others with varying numbers of handshakes. Record or graph your results.
Can you observe any patterns or make any generalizations about the experiments you just tried?
trieu :
Extension Problem: Can you find a formula for the total number of handshakes for a group of <i>n</i> people?

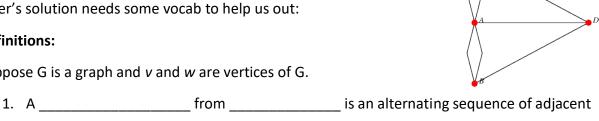
Section 3.4: The Konigsberg Bridge Problem

The beginning of this unit posed the Konigsberg Bridge Problem: "Is it possible for a person to walk around the city crossing each bridge exactly once, starting and ending at the same point?"

Euler's solution needs some vocab to help us out:

Definitions:

Suppose G is a graph and v and w are vertices of G.



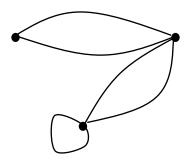
2.	Α	from	is a walk from	to	
		 			
	in which no edge is reg	eated.			

vertices and edges beginning with and ending with .

3.	Α	is a	path that starts and	l ends at the same vertex.

Using this language, we can no re-state the problem: "Does the Konigsberg Bridge problem have an Euler Circuit?"

Use this example to help you out. (Label all of the vertices and edges.)



Name a walk from v_1 to v_3 . Is that walk a path? Why or why not?

Name a circuit from v_1 to v_1 .

Is there an Euler Circuit in this graph? If so, find it. If not, explain why.

Fill in the table below to help show the difference	es among walks	, paths, circuits	s, and Euler
circuits (Fill in each box with "yes", "no", or "allo	wed")		

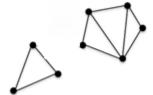
	Repeated edge?	Start/End same point?	Includes all vertices/edges?
Walk			
Path			
Circuit			
Euler			
Circuit			

Circuit				
Euler Circuit				
Euler Circui	t Theorem:			
If a graph h degree.	as an Euler circuit, t	hen every	of the grapl	n has an
Knowing th work?	e definition above,	you can now solve	the Konigsberg Bridge F	Problem Will it
D Us		C' - 'I Thanks		
			ork? Is it true that if ever c, show a graph that wo	•
This leads t	o the concept of 'co	nnectedness'.		
Definition:	Suppose G is a grap	h. Two vertices <i>v</i> a	nd w in G are	vertices if
and only if	there is a	in G fror	n v to w.	

Example 1:

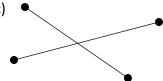
Tell whether or not each graph is connected:

a)



b)



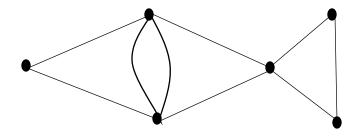


With the idea of connectedness, we can add on to the Euler circuit theorem:

Theorem: If a graph G is connected and every vertex of G has ______ degree, then G has an Euler circuit.

Example 2:

Does the following graph have an Euler circuit? If so, find such a circuit.



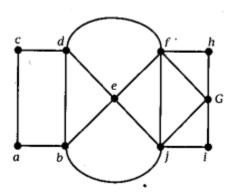
Theorem:

If a _____ and an edge is removed

from the circuit, then the resulting graph is also ______.

Section 3.5: Hamiltonian Circuits and Graph Coloring

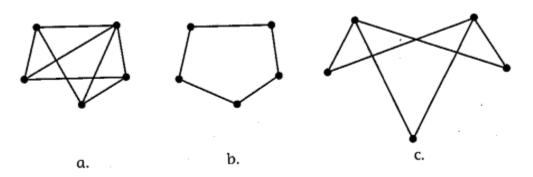
Suppose that you are a city inspector, but instead of inspecting all of the streets in an efficient manner, you must inspect the fire hydrants that are located at every intersection. This implies that you are searching for an optimal route that begins at garage, G, visits each intersection exactly once, and returns to the garage.



Find a path that meets this criteria:	
A path that uses each vertex of a graph exactly once is known as a	
If the path ends at the starting vertex, it is called a	

Example 1:

Find a Hamiltonian circuit for each graph below (if one exists):



As with Euler circuits, often it is useful for the edges of the graph to have direction. Consider a competition in which each player must play every other player. By using directed edges, it is possible to indicate winners and losers. To illustrate this, draw a complete graph in which the vertices represent the players, and a directed edge from A to B indicates that A defeats player B. This type of graph is known as a **tournament**.

Every tournament contains at least one	If there is exactly
one, it can be used to rank the teams in order, from winner to loser.	

Example 2:

Suppose four teams play in the school soccer round-robin tournament. The results of the competition follow:

Game	AB	AC	AD	ВС	BD	CD
Winner	В	Α	D	В	D	D

Draw a graph to represent the tournament. Find a Hamiltonian path and use it to rank the participants from winner to loser.

Graph Coloring

When it's time to schedule meetings or register for new classes, scheduling conflicts often arise. Mathematicians have found that graphs are useful tools in helping to resolve these conflicts.

Explore This: Below is a table of clubs at Central High School and students who hold offices in these clubs:

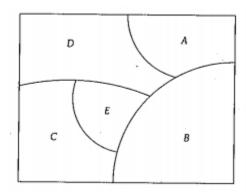
	Math Club	Chess Club	Science Club	Art Club	Pep Club	Spanish Club
Nicole	х	х	х			
Maddie	х			Х	Х	
Quinn		х				х
Emma	х		х			
Alex	х				Х	

Each club wants to meet once a week. Since several students hold offices in more than one organization, it is necessary to arrange the meeting days so that no students are scheduled for more than one meeting on the same day. Is it possible to create such a schedule? What is the minimum number of days needed (we want the fewest number of days possible)?

Problems of this type are called **coloring** problems because historically the labels placed on the vertices of the graphs were referred to as *colors*. The process of labeling the graph is called **coloring the graph**. The minimum number of labels that can be used is known as the **chromatic number** of the graph. What is the chromatic number for the graph above?

Example 3:

Use the concept of coloring to color the following map using four or fewer colors.

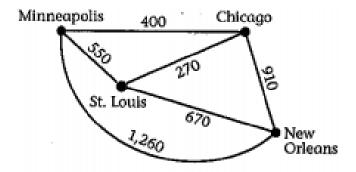


- 1. Represent the map with a graph where each vertex represents a region on the map.
- 2. Draw edges between vertices if the regions on the map have a common border.
- 3. Label the graph using a minimum number of colors.

Section 3.6: The Traveling Salesperson Problem/Finding the Shortest Route

A **traveling salesperson problem (TSP)** involves finding a Hamiltonian circuit or minimum value such as time, distance or cost. Optimization problems of this type are becoming increasingly important in the world of communications, warehousing (AMAZON!), airline networking, delivery truck routing, and building wiring.

<u>Example 1:</u> Suppose you are a salesperson who lives in St. Louis. Once a week, you have to travel to Minneapolis, Chicago, and New Orleans and then return home. The graph below represents the trips that are available to you. The edges represent cost. When each edge is assigned a number, we call it a **weighted graph**.



To save money, we want to find the least expensive route that begins in St. Louis, visits each of the other cities exactly once, and returns home. One way to solve this problem is the **brute force method**, by listing every possible circuit. Draw a tree diagram to list each possible circuit.

What if we didn't need to visit every vertex and return to the starting point, but instead you only needed to find the shortest path from one vertex in the graph to another?

Shortest Path Algorithm

- 1. Label the starting vertex S and circle it. Examine all edges that have S as an endpoint. Darken/highlight the edge with the shortest length and circle the vertex at the other endpoint of the darkened edge.
- 2. Examine all uncircled vertices that are adjacent to the circled vertices.
- 3. Using only circled vertices and darkened edges between the vertices circled, find the lengths of all paths from S to each vertex being examined. Choose the vertex and the edge that yield the shortest path. Circle this vertex and darken this edge. Ties are broken arbitrarily.
- 4. Repeat steps 2 and 3 until all vertices are circled. The darkened edges of the graph form the shortest routes from S to every other vertex in the graph.

<u>Example 2:</u> Use the shortest path algorithm to find the shortest path from A to F in the graph.

